



How to backup and pwn using Time Machine

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About me

- Nguyen Hoang Thach (@hi_im_d4rkn3ss)
- Security Researcher at STAR Labs SG Pte. Ltd.
- Focusing on Virtual Machine / Android / IOT bug hunting.
- Participated in Pwn2Own Tokyo 2020 and Pwn2wn Austin 2021 in Router, NAS and Mobile phone category, and Pwn2Own Vancouver 2022 in the Virtual Machine category

Agenda

- **1**. File service in NAS devices
- 2. Bug in processing `appl` file
- 3. Bug in processing file 's `xattr`
- 4. Summary

Nas devices

- Network-attached storage (NAS) device is a data storage device that connects to and is accessed through a network, instead of connecting directly to a computer.
- In recent year, ZDI added some NAS devices to list target (WD, Synology NAS) in their Pwn2Own contest
- Last year, I participated in Pwn2Own, I found 4 bugs and successful pwned 3 different NAS devices: WD Home Cloud NAS (release version), WD Home Cloud NAS (beta version) and WD Pro PR4100 NAS.
- Attack surface: File service

Architecture

 WD Home Cloud (release version) Arm 32bit little endian
 WD Home Cloud (beta version) Arm 32bit little endian
 WD Pro PR4100

Arch64 little endian

File Service

- Usually, NAS devices implemented at least one File Service to support file sharing, file printing, file backup.
- I will focus on 2 popular file services: *netatalk afpd* and *samba smbd*
- WD Home Cloud (release version) and WD Pro Pr4100 implement both *afpd* and *smbd*
- WD Home Cloud (beta version) implement *smbd*
- Version:
 - Netatalk afpd : v3.1.12
 - Samba smbd : v4.9.5

Configuration

• Usually, in NAS devices, at least, there is one public share folder.

• Some features also are implemented, for example: *Time Machine Backup*

-> extend the attack surface

afpd configuration

```
[ Global ]
uam list = uams_guest.so,uams_dhx2_passwd.so
save password = no
unix charset = UTF8
use sendfile = yes
zeroconf = no
guest account = nobody
vol dbpath = /data/wd/diskVolume0/backups/.systemfile/netatalk/CNID
...
[ TimeMachineBackup ]
path = /data/wd/diskVolume0/backups/timemachine
ea = auto
...
```

- `uams_guest.so` is declared in `uam list`, it accepts guest authentication.
- *TimeMachineBackup* is a public share folder

smbd configuration

- *TimeMachineBackup* is a public share folder
- `guest ok = yes` is declared, it allows guest authentication
- The `vfs objects` list contains 3 modules:
 catia, *fruit*, *streams_xattr*
- *vfs_fruit*: Enhanced OS X and Netatalk interoperability

[global]

....

[TimeMachineBackup] path = /data/wd/diskVolume0/backups/timemachine browseable = yes public = yes available = yes oplocks = yes follow symlinks = yes map archive = no guest ok = yes writable = yes vfs objects = catia fruit streams xattr durable handles = yes kernel oplocks = no kernel share modes = no posix locking = no inherit acls = yes strict sync = yes fruit:time machine = yes fruit:time machine max size = 0M

Mitigation

• afpd (WD Home)

• afpd	(WD Pro)	

• smbd	(WD Home Beta)
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	s checksec ./WDHome/afpd [*] '/tmp/WDHome/afpd'				
Arch:	arm-32-little				
RELRO:	Partial RELRO				
Stack:					
NX:	NX enabled				
PIE:	PIE enabled				
RUNPATH:	b'/system/lib:/space/jenkins/workspace/netatalk/db//rootfs/				

\$ checksec .	\$ checksec ./WDPro/afpd			
[*] '/tmp/WC	Pro/afpd'			
Arch:	amd64-64-little			
RELRO:	Partial RELRO			
Stack:				
NX:	NX enabled			
PIE:	PIE enabled			

\$ checksec ./WDHomeBeta/smbd [*] '/tmp/WDHomeBeta/smbd'				
Arch:	aarch64-64-little			
RELRO:	Full RELRO			
Stack:	Canary found			
NX:	NX enabled			
	PIE enabled			
RUNPATH:	b'/usr/lib/aarch64-linux-gnu/samba'			
FORTIFY:	Enabled			

Target

• 2 bugs in *afpd*

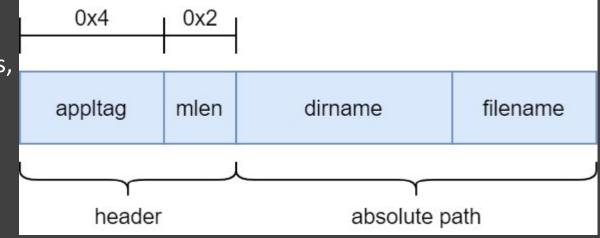
• Bugs were used to exploit the WD Home Cloud (release version)

Background

- *appl* file store database information when user read/write to files.
- In *afpd*, it has extension ".appl" and it is stored in `dbpath` which is declared in afpd.conf
- In *afpd*, there are 2 functions to create/delete *appl* file: `afp_addappl` and `afp_rmvappl`. Both functions require authentication to access. When call these functions, user will submit a `creator` value, then based on this value, a *appl* file is processed.



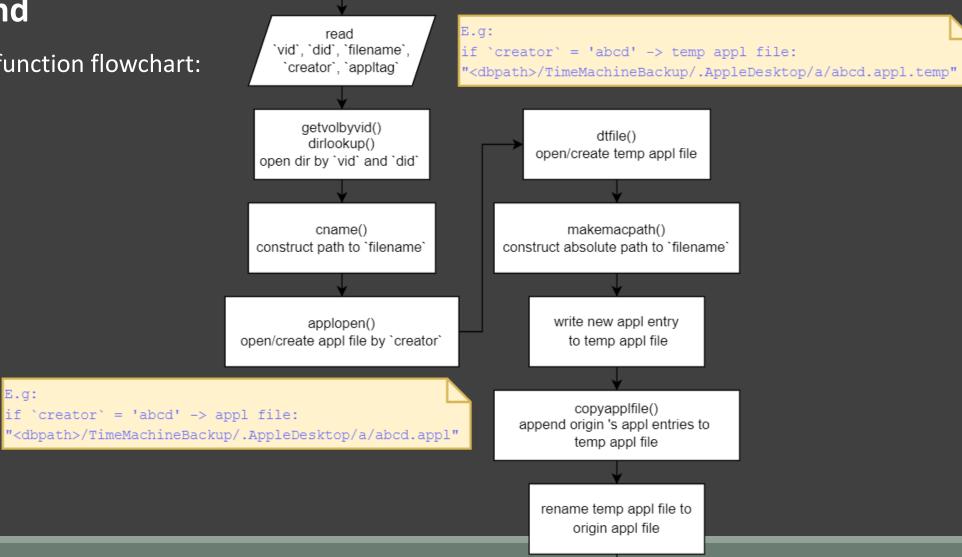
- `appltag`: 0x4 bytes, user supplied
- `mlen`: 0x2 bytes, size of absolute path
- `absolute path`: maximum 0x1000 bytes, it is absolute path of requested file by user.



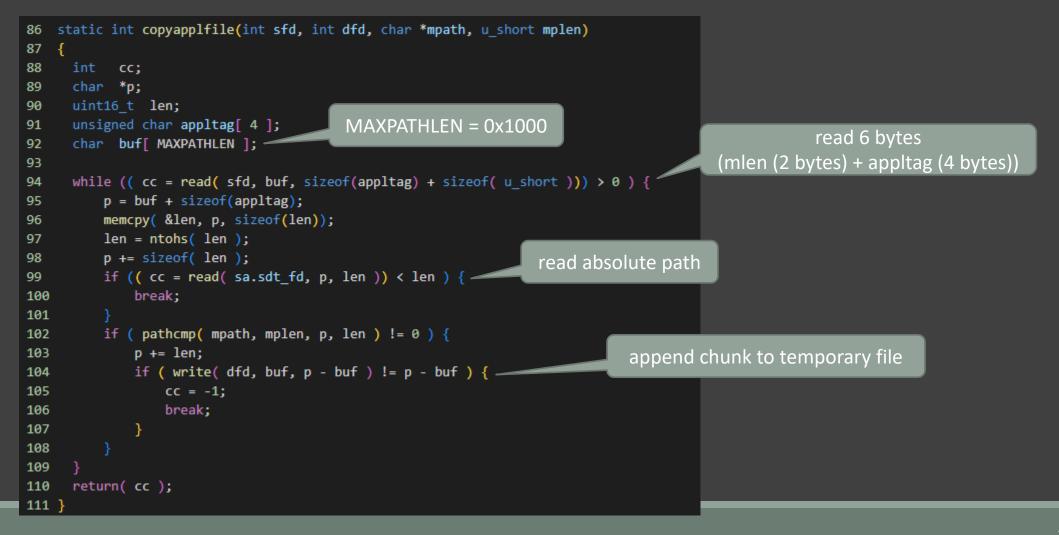
Background

`afp_addappl` function flowchart:

E.q:



Stack Out-Of-Bounds Write # root cause



Stack Out-Of-Bounds Write # root cause

- As mentioned above, maximum size of a chunk is 0x1000 + 6 = 0x1006
- `buf` is a stack-based buffer, size 0x1000
- Calling `read` function at line 99 cause Stack Out-Of-Bounds Write

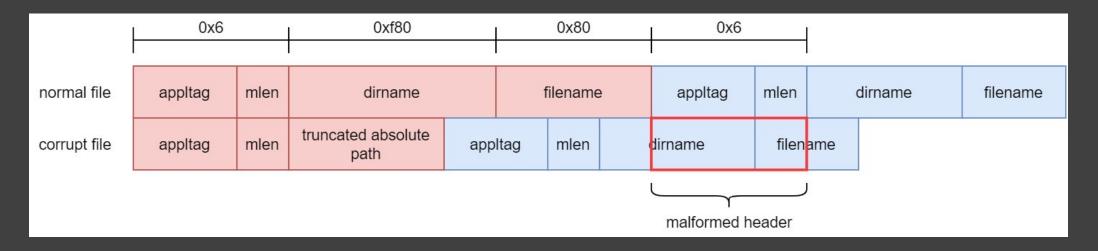
-00001026	buf	DCB	40	96	j dup(?)
-00000026		DCB	?	;	undefined
-00000025		DCB	?	;	undefined
-00000024		DCB	?	;	undefined
-00000023		DCB	?	;	undefined
-00000022	len	DCW	?		

- With 6 bytes Out-Of-Bounds Write -> overwrite `len`
- Calling `write` function at line 104 cause writing a truncated chunk to temporary *appl* file

Stack Out-Of-Bounds Write # root cause

• The next time function `copyapplfile` parses corrupted *appl* file, the calling `read` at line 99 might cause Stack Out-Of-Bounds Write, and we could overwrite return address in stack -> RCE

• Here is a sample payload cause corrupting *appl* file:



Stack Out-Of-Bounds Write # exploitation

- This bug is used to exploit WD Home Cloud (release version)
- Architecture: arm 32 bit
- Mitigation: ASLR + PIE

Notes

- The maximum size of filename in linux is 256, we need 0x1000 bytes -> we need to create multiple nested folder.
- Since absolute path cannot contains null char -> cannot store pointer address in it
- Red filename will overwrite the `len` value in stack, I set it to 0xf60, when translate to ascii, it is `\x60\0xf`, still valid to use in filename
- sizeof `appItag` == 4 bytes and controllable by user -> we will place malform `mlen` and malform return address in it.

Stack Out-Of-Bounds Write # exploitation

Step 1. Bypass ASLR

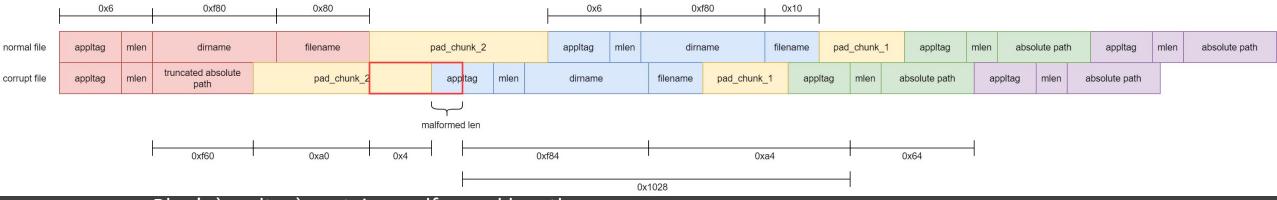
- *afpd* is multi-process server, using `fork` to create child process to handle a new connection
 -> we could partial overwrite ret address to bruteforce PIE base address.
- Partial overwrite origin ret address not work, because \$r11 register is overwritten in stack and parent function use \$r11 -> always crash
- I used timebase bruteforce method instead
- 2 address is different, but difference is not too large (< 0x2000) -> bruteforce still work
- .text:0000D540 BL copyapplfile .text:0000D544 STR R0, [R11,#var_24] .text:0000ECC8 MOVW R0, #2 ; seconds .text:0000ECCC BL sleep

- 1st byte is always in range 0xa0 0xaf
- -> maximum 16 + 256 + 16 = 288 attempts to successful bruteforce PIE base address

Stack Out-Of-Bounds Write # exploitation

Step 2. Execute command	as root			`afprun_bg` function run
	.text:000268A4	LDR	R1, [SP,#0x64]	command with root privilege
• Use the following gadget:	.text:000268A8	BL	afprun_bg	

- Since the file content is copied to stack -> we could put address of command in file content
- The final payload:



- Blue's `appItag` contains malformed length
- Green's `appltag` contains address of above gadget
- Purple's `appltag` contains address of command

Race condition # root cause

```
201 int afp addappl(AFPObj *obj, char *ibuf, size t ibuflen U, char *rbuf U, size t *rbuflen)
202 {
                                                                                   Open/Create the
       if (( tfd = open( tempfile, 0 RDWR|0 CREAT, 0666 )) < 0 ) { -</pre>
251
                                                                                    `tempfile` to edit
252
           return( AFPERR PARAM );
253
       mpath = obj->newtmp;
254
       mp = makemacpath( vol, mpath, AFPOBJ TMPSIZ, curdir, path->m name );
255
256
       if (!mp) {
           close(tfd);
257
258
           return AFPERR PARAM;
259
       mplen = mpath + AFPOBJ_TMPSIZ - mp;
260
261
       /* write the new appl entry at start of temporary file */
262
       p = mp - sizeof( u short );
263
       mplen = htons( mplen );
264
                                                          Contruct new chunk
       memcpy( p, &mplen, sizeof( mplen ));
265
266
       mplen = ntohs( mplen );
       p -= sizeof( appltag );
267
268
       memcpy(p, appltag, sizeof( appltag ));
                                                         Append new chunk to
269
       cc = mpath + AFPOBJ_TMPSIZ - p;
                                                                `tempfile`
        if ( write( tfd, p, cc ) != cc ) { ----
270
271
           close(tfd);
272
           unlink( tempfile );
           return( AFPERR_PARAM );
273
274
288
```

Race condition # root cause

- *afpd* is a multiple processes service each command is processed in a separated process
- At line 251 and 270, perform file operator without lock.
- Sending multiple add appl file commands with same `creator` value -> multiples process
 processed a same file -> race condition
- Race condition -> chunks data might overlap each other -> corrupt the temporary appl file

Race condition # root cause

 As mentioned before, when function `copyapplfile` parses corrupted *appl* file, the calling `read` at line 99 might cause Stack Out-Of-Bounds Write, and we could overwrite return address in stack -> RCE

• Here is a sample payload cause corrupting *appl* file:

process 1 's chunk	appltag_1	mlen_1	dirname_1			filename_1
process 2 's chunk	appltag_2	mlen_2	dirname_2	filename_2		
corrupt file	appltag_2	mlen_2	dirname_2	filename_2		filename_1
					malfo	rmed header

Race condition # exploitation

This bug is used to exploit WD Home Cloud (release version)

Architecture: arm 32 bit

Mitigation: ASLR + PIE

Race condition # exploitation

Step 1. Bypass ASLR

- Can reuse timebased bruteforce ?
- Race condition + bruteforce seems not reliable
- Need a information disclosure vulnerability

Race condition # exploitation

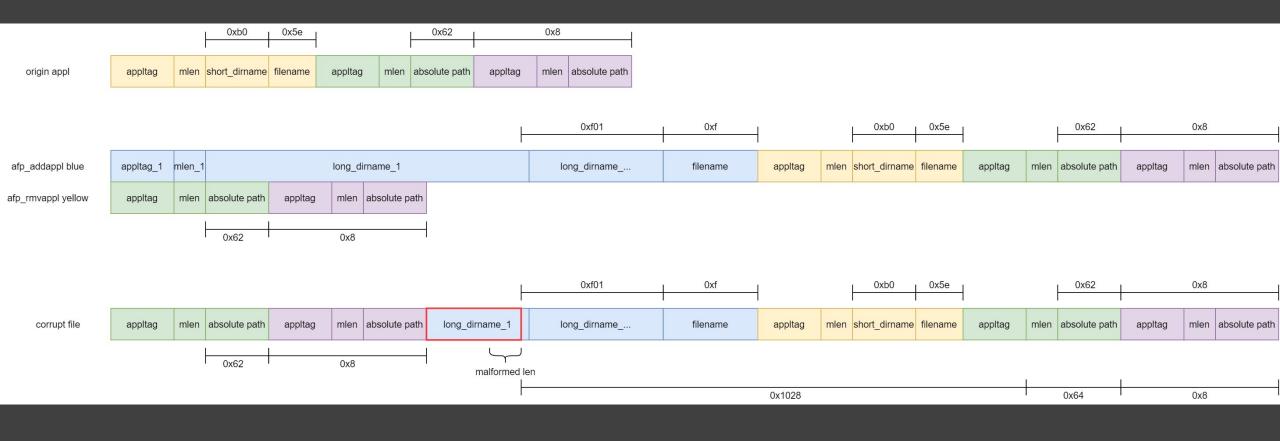
Step 2. Execute command as root

- Use the following gadget: .text:000268A4 LDR R1, [SP,#0x64] command with root privilege afprun_bg
- Since the file content is copied to stack -> we could put address of command in file content
- The race condition also occurred in `afp_rmappl` -> could race between `afp_addappl` and `afp_rmappl` processes

`afprun bg` function run

Race condition # exploitation

Step 2. Execute command as root



Race condition # exploitation

- Step 2. Execute command as root
- Malformed length is lied on the end of Blue's directory name
- I choose `Malformed length` is 0x1108, translated to ascii name: `\x08\x11` still valid to use as directory name
- Green's `appltag` contains address of gadget
- Purple's `appltag` contains address of command

Bonus

Netatalk weak hash function leads to information disclosure

- Bug is in "uams_dhx2_passwd.so"
- Using weak hash function to hash a pointer ??



• *Session ID* value here will be sent back to client later

Bonus

Netatalk weak hash function leads to information disclosure

- `obj` is a global pointer -> located in the .text section
- The NAS running 32bit OS -> 1st byte and 4th byte are known
- We could calculate the 2nd byte and 3rd byte from *Session ID*
 - -> bypass ASLR

Target

• 2 bugs: one in *afpd* and one in *smbd*

• Bugs were used to exploit the WD Pro PR4100 and WD Home Cloud (beta version)

Background

- Extended attributes (xattr) are *name:value* pairs associated permanently with files and directories
- Both *afpd* and *smbd* have command to allow user to set xattr for a file/directory (require authentication).
- Some special xattr will be parsed when process files

Background

In case of *afpd*:

 `afp_setextattr` command is responsible to set the *value* of the extended attribute identified by *name* and associated with the given path in the filesystem.

• It is done by invoking `setxattr`/`lsetxattr`/`fsetxattr` function.

• No checking in whole process -> user can set arbitrary *name*:*value* xattr

Background

• `ad_open` function is responsible to open file.

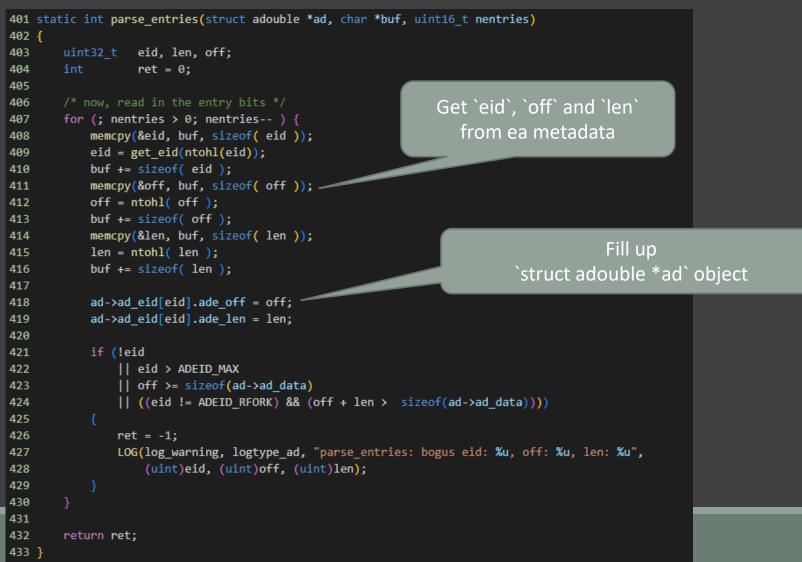
- Some special xattrs are parsed here, one of them is AD_EA_META: "org.netatalk.Metadata"
- As mentioned before, no checking in the `afp_setextattr` function -> user can set the malform "org.netatalk.Metadata" xattr.

Background

`parse_entries` function call stack



Afpd Parsing xattr Out-Of-Bounds Access # root cause



Afpd Parsing xattr Out-Of-Bounds Access # root cause

• Each entry in `struct adouble *ad` object has format:

0x4	0x4	0x4		
eid	off	len		

`eid`: ID of entry

- `off` is offset value from `ad->ad_data` buffer
- `len` is size of value.

 There are some checks to prevent accessing out-of-bounds of `ad->ad_eid` array and `ad->ad_data` buffer

 But no check if the `off` and `len` are valid for a specific `eid` -> leads to multiple Out-Of-Bounds access when use `struct adouble *ad` later.

Afpd Parsing xattr Out-Of-Bounds Read # root cause

```
void *get finderinfo(const struct vol *vol, const char *upath, struct adouble *adp, void *data, int islink)
80
81
82
       struct extmap
                           *em;
83
                           *ad finder = NULL;
84
                           chk ext = 0;
85
                                                            (adp->ad data + adp->ad eid[ADEID FINDERI].ade off)
86
       if (adp)
87
           ad finder = ad entry(adp, ADEID FINDERI);
88
       if (ad_finder) {
89
           memcpy(data, ad finder, ADEDLEN FINDERI); // <-- ADEDLEN FINDERI = 32
90
           /* default type ? */
91
           if (default type(ad finder))
92
93
               chk ext = 1;
94
       else {
95
122
```

Afpd Parsing xattr Out-Of-Bounds Read # root cause

- Line 87, get `ad_finder` pointer from `struct adouble *adp`. As mentioned above, `ad_finder` could point to the last byte of `adp->ad_data` buffer.
- Line 90, calling `memcpy` with the fixed size 32, lack of the check if 32 > `adp->ad_eid[ADEID_FINDERI].ade_len` -> out-of-bounds read issue.
- The `data` will be sent back to user later -> information disclosure
- `adp->ad_data` is a stack-based buffer -> leak pie base address.

Afpd Parsing xattr Out-Of-Bounds Write # root cause

```
833 static int ad addcomment(const AFPObj *obj, struct vol *vol, struct path *path, char *ibuf)
834 {
835
       struct ofork
                           *of:
                           *name, *upath;
836
837
                           isadir:
838
               clen;
        int
       struct adouble ad, *adp;
839
840
                                           Get `clen` from user-supplied buffer
841
       clen = (u char)*ibuf++;
       clen = min( clen, 199 );
842
       if (ad getentryoff(adp, ADEID COMMENT)) {
862
863
           if ( (ad_get_MD_flags( adp ) & O_CREAT) ) {
               if ( *path->m name == '0' ) {
864
                   name = (char *)curdir->d m name->data;
865
               } else {
866
867
                   name = path->m name;
868
869
               ad setname(adp, name);
870
                                                                            (adp->ad data + adp->ad eid[ADEID COMMENT].ade off)
           ad_setentrylen( adp, ADEID_COMMENT, clen );
871
           memcpy( ad entry( adp, ADEID_COMMENT ), ibuf, clen );
872
873
           ad flush( adp );
874
875
       ad_close(adp, ADFLAGS_HF);
876
       return( AFP_OK );
877 }
```

Afpd Parsing xattr Out-Of-Bounds Write # root cause

- Line 872, get pointer from `struct adouble *adp`. As mentioned above, this pointer could point to the last byte of `adp->ad_data` buffer.
- Line 872, calling `memcpy` with the controllable len `clen`, lack of the check if `clen` > `adp->ad_eid[ADEID_FINDERI].ade_len` -> out-of-bounds write issue.
- `adp->ad_data` is a stack-based buffer -> could overwrite the return address in stack.

Afpd Parsing xattr Out-Of-Bounds Access # exploitation

This bug is used to exploit WD Pro Pr4100

Architecture: aarch64

Mitigation: ASLR + PIE

Afpd Parsing xattr Out-Of-Bounds Access # exploitation

Step 1: Bypass ASLR

• Using Out-Of-Bounds Read to leak PIE base

Step 2: Execute command as root

 Using Out-Of-Bounds Write to overwrite the return address in stack with the following rop chain:

0x000000000003c429 : pop rsi ; pop r15 ; ret

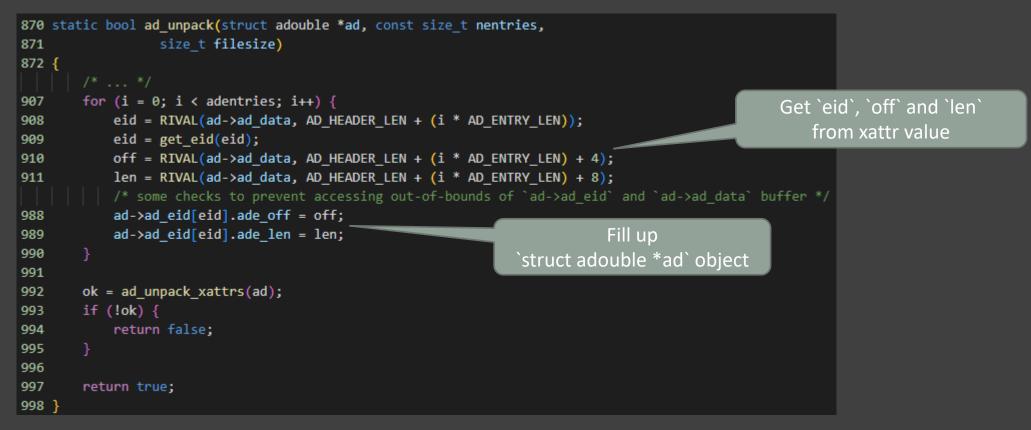
Follow by address of a global buffer stored command + address of `afprun_bg` function

-> execute command as root

Smbd Parsing xattr Out-Of-Bounds Access # root cause

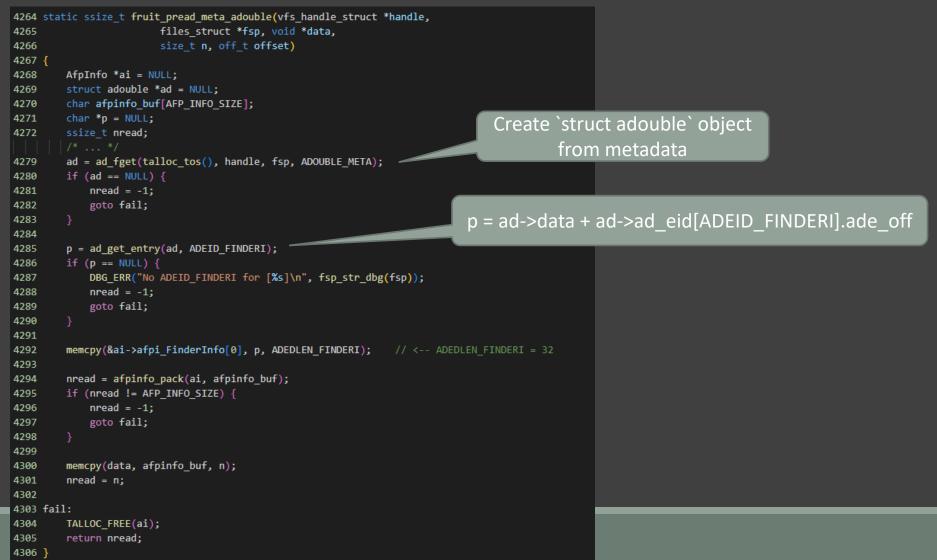
- *smbd* also provide to user a command to set xattr of a file/directory.
- It is done by `set_ea` function.
- The *name* of our attribute must not be in the private *Samba* attribute name list (`user.SAMBA_PAI`, `user.DOSATTRIB`, `user.SAMBA_STREAMS`, `security.NTACL`)
- When *fruit* module process file/directory, it also parse some special xattr values, such as AFPINFO_EA_NETATALK: "org.netatalk.Metadata"
- `AFPINFO_EA_NETATALK` is not in the private attribute name list -> user can submit a malform xattr value

Smbd Parsing xattr Out-Of-Bounds Access # root cause



It is very similar to the *afpd* 's `parse_entries` function

Smbd Parsing xattr Out-Of-Bounds Read # root cause



Smbd Parsing xattr Out-Of-Bounds Read # root cause

- Line 4285, get `p` pointer from `struct adouble *ad`. As mentioned above, `p` could point to the last byte of `ad->ad_data` buffer.
- Line 4293, calling `memcpy` with the fixed size 32, lack of the check if 32 > `ad->ad_eid[ADEID_FINDERI].ade_len` -> out-of-bounds read issue.
- The `data` will be sent back to user later -> information disclosure
- `ad->ad_data` is a heap-based buffer.

Smbd Parsing xattr Out-Of-Bounds Write # root cause

```
4642 static ssize t fruit pwrite meta netatalk(vfs_handle_struct *handle,
4643
                          files struct *fsp, const void *data,
4644
                          size t n, off t offset)
4645
4646
         struct adouble *ad = NULL;
         AfpInfo *ai = NULL;
4647
4648
         char *p = NULL;
         int ret;
4649
                                                                           Create `struct adouble` object
         bool ok:
4650
                                                                                   from metadata
         ad = ad_fget(talloc_tos(), handle, fsp, ADOUBLE_META);
4657
4658
         if (ad == NULL) {
            ad = ad init(talloc tos(), handle, ADOUBLE META);
4659
            if (ad == NULL) {
4660
4661
                return -1;
4662
                                                                         p = ad->data + ad->ad eid[ADEID FINDERI].ade off
4663
4664
         p = ad_get_entry(ad, ADEID_FINDERI);
         if (p == NULL) {
4665
4666
            DBG_ERR("No ADEID_FINDERI for [%s]\n", fsp_str_dbg(fsp));
4667
            TALLOC FREE(ad);
4668
            return -1;
4669
4670
4671
         memcpy(p, &ai->afpi FinderInfo[0], ADEDLEN FINDERI);
                                                                   // <-- ADEDLEN FINDERI = 32</pre>
4672
                                                                                                                                    47
4704
```

Smbd Parsing xattr Out-Of-Bounds Write # root cause

- Line 4664, get `p` pointer from `struct adouble *ad`. `p` pointer could point to the last byte of `ad->ad_data` buffer.
- Line 4671, calling `memcpy` with the fixed size 32, lack of the check if 32 > `ad->ad_eid[ADEID_FINDERI].ade_len` -> out-of-bounds write issue.
- `ad->ad_data` is a heap-based buffer.

Smbd Parsing xattr Out-Of-Bounds Access # exploitation

This bug is used to exploit WD Home Cloud (beta version)

- Architecture: arm 32 bit
- Mitigation: ASLR + PIE

Smbd Parsing xattr Out-Of-Bounds Access # exploitation

Backgound

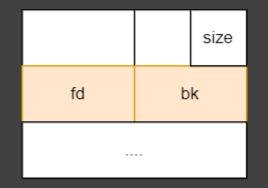
• *Smbd* implements both glibc 's `ptmalloc` and its own `talloc` memory allocation.

• Chunk format:

flags	next
prev	parent
child	refs
destructor	name

Talloc chunk

- flags: chunk canary + some flags
- next, prev: point to the next/prev chunk
- destructor: function pointer, will be invoked when chunk is free



Ptmalloc chunk:

• fd, bk: point to the next/prev chunk

Smbd Parsing xattr Out-Of-Bounds Access # exploitation

Step 1. Bypass ASLR

- Require spray heap to create both `ptmalloc` chunk and `talloc` chunk
- Able to read up to 24 bytes past end of `ad->ad_data` chunk
 - Leak `talloc.flags` to bypass chunk canary check
 - Leak `talloc.next`, `talloc.prev` to know heap address
 - Leak `ptmalloc.fd`, `ptmalloc.bk` to know libc address (because it might point to main_arena)

Smbd Parsing xattr Out-Of-Bounds Access # exploitation

Step 2. Control the \$pc

• We know chunk canary -> can forge a valid `talloc` chunk in heap address

 We know heap address -> overwrite `next` pointer in a in used `talloc` chunk by our forge chunk -> when this in used `talloc` chunk is freed -> our forge chunk is also freed -> invoke `destructor` function

 We know libc address -> calculate other shared lib address -> set proper address in `destructor` function pointer -> execute command as root

Summary

Conclusion

 Check configuration of file service running on NAS/router devices, it might contains addition feature -> extend attack surface

• Same feature might contain same bug pattern

TODO

- Does the fuzzer work ?
- Samba: Check remain modules which declared in `vfs objects` list





Thanks for listening

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